THE GEORGE WASHINGTON UNIVERSITY

WASHINGTON, DC

5b. Normalization

CSCI 2541 Database Systems & Team Projects

Wood & Chaufournier

Slides adapted from Prof. Bhagi Narahari; and Silberschatz, Korth, and Sudarshan

Announcements?

Exam! ~1.5 – 2.5 weeks away

Relational Algebra HW ✓

SQL and Shopping Cart HW

- ER diagrams
- SQL

Last time…

SQL DDL & DML

Entity Relationship Model

Normalization
- Bad Schemas
- Normal Forms
- Functional Dependencies

this time…

Good Schemas

The ER model can help us design a logical DB structure that matches our business goals

The conceptual schema must be translated into a logical (SQL) schema

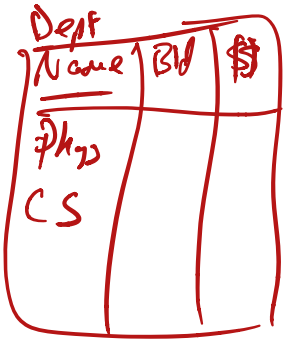
How do we judge if a SQL schema is well designed?

Bad Schemas

Let's track professors and their department

- We will put all the info together in one table so we don't have to worry about joining stuff!

ID	*name*	*salary*	*dept_name*	*building*	*budget*
22222	Einstein	95000	Physics	Watson	70000
12121	Wu	90000	Finance	Painter	120000
32343	El Said	60000	History	Painter	50000
45565	Katz	75000	Comp. Sci.	Taylor	100000
98345	Kim	80000	Elec. Eng.	Taylor	85000
76766	Crick	72000	Biology	Watson	90000
10101	Srinivasan	65000	Comp. Sci.	Taylor	100000
58583	Califieri	62000	History	Painter	50000
83821	Brandt	92000	Comp. Sci.	Taylor	100000
15151	Mozart	40000	Music	Packard	80000
33456	Gold	87000	Physics	Watson	70000
76543	Singh	80000	Finance	Painter	120000



Why is this a bad idea?

Bad Schemas

Let's track professors and their department

- We will put all the info together in one table so we don't have to worry about joining stuff!

ID	*name*	*salary*	*dept_name*	*building*	*budget*
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58583	Califieri	62000	History	Painter	50000

Update Anomalies: need to modify all repetitive rows

Insertion Anomalies: Need to use NULL if we add a department with no instructors

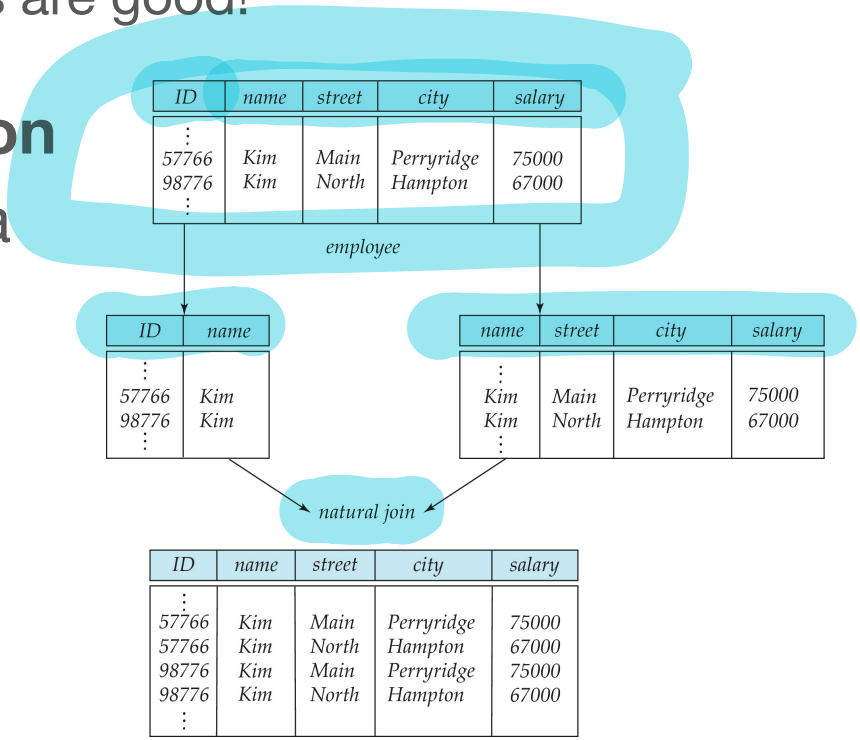
Deletion Anomalies: Removing all instructors loses information about the department

Splitting Tables

Decomposing into separate tables helps resolve this… but there are multiple ways to split tables

- Not all decompositions are good!

A **Lossy Decomposition** results in us losing data if we try to merge back using a join



What is Normalization?

1. Tests to see how “good” a schema is

2. Normalization algorithms to decompose relations into smaller relations that contain less redundancy
   - This decomposition requires that **no information is lost** and **reconstruction** of the original relations from the smaller relations must be possible.

Normalization should be done when you design your schema and anytime you update it

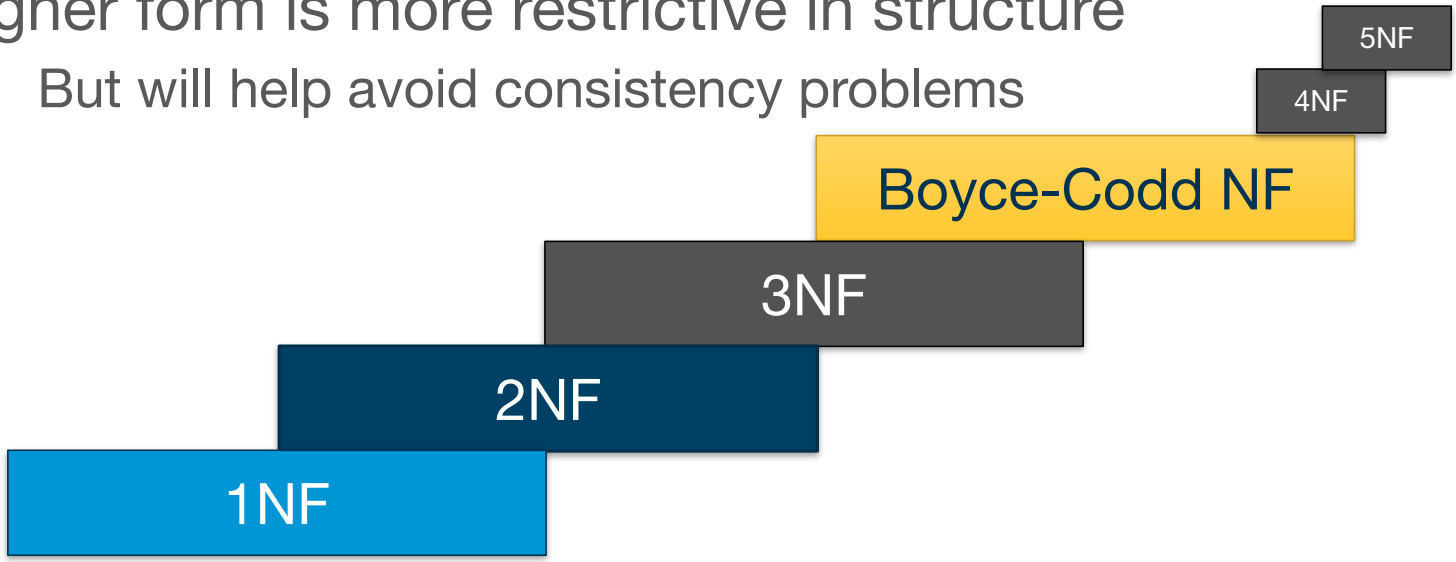
Normal Forms

Normal forms give us a hierarchy of rules

- No normalization - unconstrained, messy data
- First Normal Form - removes some redundancy
- Second Normal From - removes more redundancy… etc

Higher form is more restrictive in structure

- But will help avoid consistency problems



First Normal Form (1NF)

Attributes should be atomic and tables should have no repeating groups

Each field only has one value

No columns repeat the same type of information

No duplicate rows in the table; order doesn't matter

1NF Examples

Attributes should be atomic and tables should have no repeating groups

Do these violate 1NF?

1

Customer ID	First Name	Surname	Telephone Number
123	Pooja	Singh	555-861-2025, 192-122-1111
456	San	Zhang	(555) 403-1659 Ext. 53, 182-929-2929
789	John	Doe	555-808-9633

Not atomic!

2

Customer ID	First Name	Surname	TNumber1	TNumber2
123	Pooja	Singh	555-861-2025	192-122-1111
456	San	Zhang	(555) 403-1659 Ext. 53	182-929-2929
789	John	Doe	555-808-9633	

TN3

Examples from https://en.wikipedia.org/wiki/First_normal_form

1NF Examples

Attributes should be atomic and tables should have no repeating groups

Do these violate 1NF?

Customer ID	First Name	Surname	Telephone Number
123	Pooja	Singh	555-861-2025, 192-122-1111
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789	John	Doe	555-808-9633

Both are bad!

Customer ID	First Name	Surname	TNumber1	TNumber2
123	Pooja	Singh	555-861-2025	192-122-1111
456	San	Zhang	(555) 403-1659 Ext. 53	182-929-2929
789	John	Doe	555-808-9633	

Examples from https://en.wikipedia.org/wiki/First_normal_form

1NF Split or Flatten

Attributes should be atomic and tables should have no repeating ~~groups~~ columns

Possible solutions

Customer ID	First Name	Surname	Telephone Number
123	Pooja	Singh	555-861-2025
123	Pooja	Singh	192-122-1111
456	San	Zhang	182-929-2929
456	San	Zhang	(555) 403-1659 Ext. 53
789	John	Doe	555-808-9633

meets 1NF

OR

Customer ID	First Name	Surname
123	Pooja	Singh
456	San	Zhang
789	John	Doe

Customer ID	Telephone Number
123	555-861-2025
123	192-122-1111
456	(555) 403-1659 Ext. 53
456	182-929-2929
789	555-808-9633

Examples from https://en.wikipedia.org/wiki/First_normal_form

1NF Violations

Generally easy to detect:

1. Check for Column names with a number (telephone1, telephone2, course1, course2, etc)

2. Make sure that order of rows doesn't matter

3. Have a primary key to enforce uniqueness across rows

Second Normal Form (2NF)

> No value in a table should be dependent on only **part** of a key that uniquely identifies a row

It must be in 1NF and…

We should **not** be able to derive the value of a column based on only **a part of a Candidate Keys**

- Must hold for all Candidate Keys if there are multiple

Reminder: Key types

Superkey of R:

- A (**possibly larger than necessary**) set of attributes that is sufficient to uniquely identify each tuple in r(R)

Candidate Key of R: A “minimal” superkey

- A **minimal set** of attributes to denote uniqueness!
- A Candidate Key is a Superkey but opposite may not be true

Primary Key: A specific Candidate Key chosen to represent a relation/table

2NF Examples

No value in a table should be dependent on only part of a key that uniquely identifies a row

Does this violate 2NF?

Customer ID	First Name	Surname	Telephone Number
123	Pooja	Singh	555-861-2025
123	Pooja	Singh	192-122-1111
456	San	Zhang	182-929-2929
456	San	Zhang	(555) 403-1659 Ext. 53
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2NF Examples

No value in a table should be dependent on only part of a key that uniquely identifies a row

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Yes!

- Our Key is (Customer ID, Telephone Number), but from Customer ID alone we could uniquely identify the name
- We could make func(CustomerID) -> (First Name, Surname)

In general, better to use the splitting method for 1NF

2NF vs 1NF

Why do we care??

1NF

Customer ID	First Name	Surname	Telephone Number
123	Pooja	Singh	555-861-2025
123	Pooja	Singh	192-122-1111
456	San	Zhang	182-929-2929
456	San	Zhang	(555) 403-1659 Ext. 53
789	John	Zhang	555-808-9633

VS

2NF

Customer ID	First Name	Surname
123	Pooja	Singh
456	San	Zhang
789	John	Zhang

Customer ID	Telephone Number
123	555-861-2025
123	192-122-1111
456	(555) 403-1659 Ext. 53
456	182-929-2929
789	555-808-9633

2NF vs 1NF

Redundant data can lead to inconsistencies if it is only partially updated!

1NF

Customer ID	First Name	Surname	Telephone Number
123	Pooja	Singh	555-861-2025
123	Pooja	Sing	192-122-1111
456	San	Zhang	182-929-2929
456	San	Zhang	(555) 403-1659 Ext. 53
789	John	Zhang	555-808-9633

VS

2NF

Customer ID	First Name	Surname
123	Pooja	Singh
456	San	Zhang
789	John	Zhang

Customer ID	Telephone Number
123	555-861-2025
123	192-122-1111
456	(555) 403-1659 Ext. 53
456	182-929-2929
789	555-808-9633

More 2NF Examples

Manufacturer	Model	Price	Manufacturer country
Forte	X-Prime	50	Italy
Forte	Ultraclean	50	Italy
Dent-o-Fresh	EZbrush	65	USA
Brushmaster	SuperBrush	34	USA
Kobayashi	ST-60	22	Japan
Hoch	Toothmaster	18	Germany
Hoch	X-Prime	50	Germany

More 2NF Examples

This avoids **Update Anomalies**

- Previously we would have had to scan all tuples if a manufacturer moved to a different country to ensure consistency

Manufacturer	Model	Price
Forte	X-Prime	45
Forte	Ultraclean	50
Dent-o-Fresh	EZbrush	65
Brushmaster	SuperBrush	34
Kobayashi	ST-60	22
Hoch	Toothmaster	18
Hoch	X-Prime	22

Manufacturer	Country
Forte	Italy
Dent-o-Fresh	USA
Brushmaster	USA
Kobayashi	Japan
Hoch	Germany

Third Normal Form (3NF)

No value should be able to be derived based on another non-key field

It must be in 2NF and…

all **non-prime attributes** depend only on the **candidate keys** and do not have a **transitive dependency** on another key

3NF Intuition

No value should be able to be derived based on another non-key field

What is the redundant information in this table?

Customer ID	First Name	Surname	Birthday	Age	Fav Color
123	Pooja	Singh	1/4/1984	37	Blue
456	San	Zhang	3/15/2001	19	Blue
789	John	Zhang	11/12/2006	14	Buff

3NF Intuition

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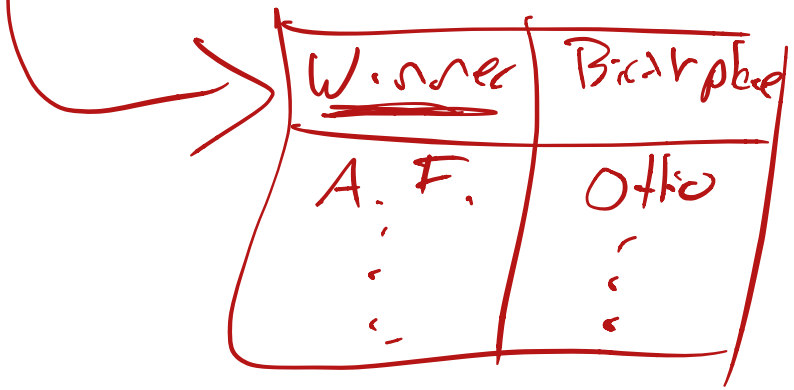
If we know Birthday, we can calculate Age -> there is an obvious dependency between them! Can remove Age.

3NF Intuition

No value should be able to be derived based on another non-key field

What is the redundant information in this table?

Tournament	Year	Winner	Winner's Birthplace
Indiana Invitational	1998	Al Fredrickson	Ohio
Cleveland Open	1999	Bob Albertson	New York
Des Moines Masters	1999	Al Fredrickson	Ohio
Indiana Invitational	1999	Chip Masterson	Kentucky



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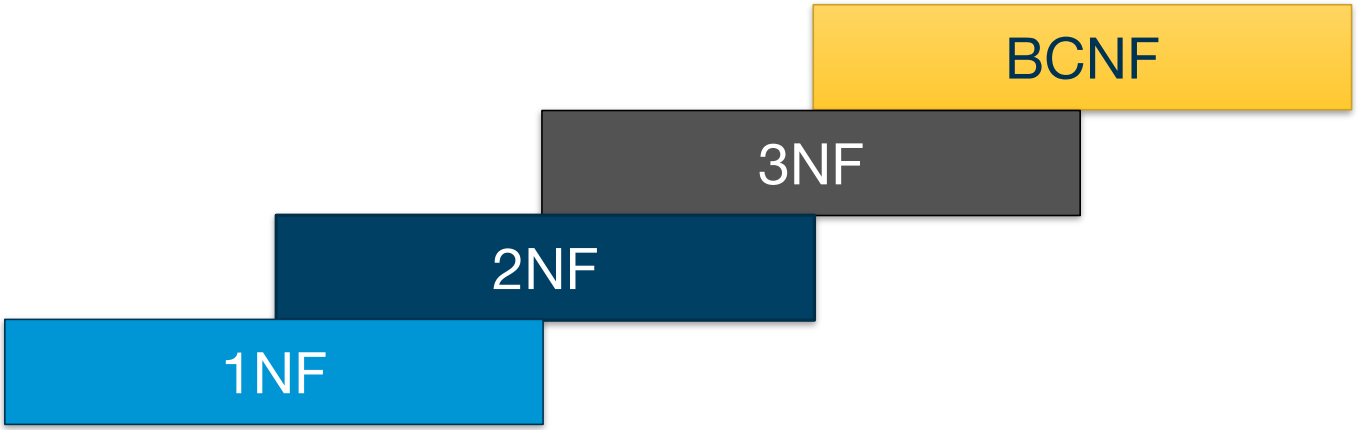
Indiana Invitational	1998	Al Fredrickson	Ohio
Cleveland Open	1999	Bob Albertson	New York
Des Moines Masters	1999	Al Fredrickson	New Jersey
Indiana Invitational	1999	Chip Masterson	Kentucky

The {Winner's Birthplace} attribute can be determined based on Winner, which is not a Candidate Key for the table. Need to split!

Normal Form Redundancy

1NF and 2NF - eliminate redundancy **across** rows

3NF, BCNF - also eliminate redundancy **within** rows



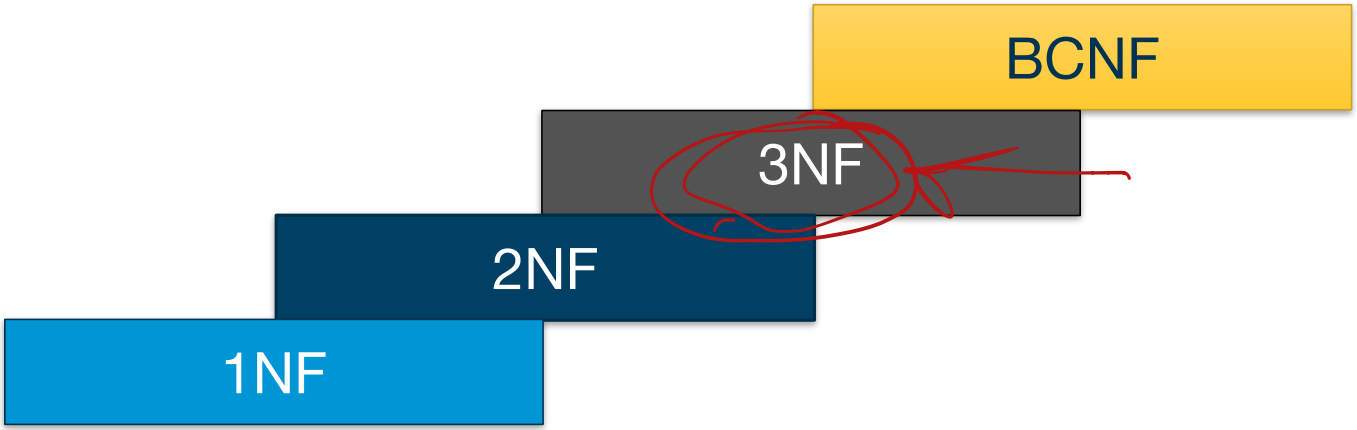
Good and Bad Schemas

Functional Dependencies

Normal Forms based on Functional Dependencies

Dependencies

How can we **formally represent dependencies** between Attributes in a Relation?



Functional Dependencies

Use **functional dependencies!** (abbreviated **FD**)

We say a set of attributes **X** functionally determines an attribute **Y** if *given the values of **X** we always know the only possible value of **Y***.

- Notation: **X** → **Y**
- X **functionally determines** Y
- Y is **functionally dependent** on X

Example:

- GWID → Name
- {GWID, CourseID, Semester, Year} → Grade , Prof Name

Sets of Functional Dependencies

Some more functional dependencies

- {GWID} → {NAME, ADDRESS, MAJOR}
- {MAJOR} → {DEPT_NAME, DEPT_CHAIR}

From above dependencies, we can infer

- {GWID} → {DEPT_NAME, DEPT_CHAIR}

We can do math on functional dependencies!

A functional dependency "holds" if it must be true for all legal relations

Functional Dependency Ops

Armstrong's Axioms: where A, B, C are sets of attributes

- **Reflexive rule**: if B ⊆ A, then A → B (if B is subset of A)
- **Augmentation rule**: if A → B, then C A → C B
- **Transitivity rule**: if A → B, and B → C, then A → C

These rules are

- Sound and complete — generate all functional dependencies that hold.

{GWID} → {Name, Address, Major}

{Major} → {Dept_Name, Dept_Chair}

{GWID, CourseID, Semester, Year} → Grade

Functional Dependency Ops

Armstrong's Axioms: where A, B, C are sets of attributes

- **Reflexive rule**: if $B \subseteq A$, then $A \rightarrow B$ (if B is subset of A)
- **Augmentation rule**: if $A \rightarrow B$, then $C\ A \rightarrow C\ B$
- **Transitivity rule**: if $A \rightarrow B$, and $B \rightarrow C$, then $A \rightarrow C$

These rules are

- Sound and complete — generate all functional dependencies that hold.

Bonus rules to make life easier:

- **Union rule**: If $\alpha \rightarrow \beta$ holds and $\alpha \rightarrow \gamma$ holds, then $\alpha \rightarrow \beta\ \gamma$ holds.
- **Decomposition rule**: If $\alpha \rightarrow \beta\ \gamma$ holds, then $\alpha \rightarrow \beta$ holds and $\alpha \rightarrow \gamma$ holds.
- **Pseudotransitivity rule**:If $\alpha \rightarrow \beta$ holds and $\gamma\ \beta \rightarrow \delta$ holds, then $\alpha\ \gamma \rightarrow \delta$ holds.

Definition: Closure of a Set of FD's

Defn. Let F be a set of FD's.

Its **closure**, **F+**, is the set of all FD's:

{X → Y | X → Y is derivable from F by Armstrong's Axioms}

Two sets of dependencies F and G are equivalent if F+=G+

- i.e., their closures are equal
- i.e., the same sets of FDs can be inferred from each

Example Closure

What FDs can we infer?

R = (A, B, C, G, H, I)

F = {A → B
A → C
CG → H
CG → I
B → H}

Reflexive rule: if $\beta \subseteq \alpha$, then $\alpha \rightarrow \beta$

Augmentation rule: if $\alpha \rightarrow \beta$, then $\gamma\,\alpha \rightarrow \gamma\,\beta$

Transitivity rule: if $\alpha \rightarrow \beta$, and $\beta \rightarrow \gamma$, then $\alpha \rightarrow \gamma$

F+

A → B
B → H
⇒ A → H

A → C
AG → CG
CG → I
⇒ AG → I

Example Closure

R = (A, B, C, G, H, I)

F = { A → B
A → C
CG → H
CG → I
B → H}

Reflexive rule: if $\beta \subseteq \alpha$, then $\alpha \rightarrow \beta$

Augmentation rule: if $\alpha \rightarrow \beta$, then $\gamma\ \alpha \rightarrow \gamma\ \beta$

Transitivity rule: if $\alpha \rightarrow \beta$, and $\beta \rightarrow \gamma$, then $\alpha \rightarrow \gamma$

A few members of F+ include:

- A → H
  - by transitivity from A → B and B → H
- AG → I
  - by augmenting A → C with G, to get AG → CG and then transitivity with CG → I
- CG → HI
  - by augmenting CG → I to infer CG → CGI, and augmenting of CG → H to infer CGI → HI, and then transitivity

Good and Bad Schemas

Functional Dependencies

Normal Forms based on Functional Dependencies

Redefining 2NF

Using Functional Dependencies and Closures lets us more precisely define our Normal Forms

Second Normal Form: For every $\mathbf{X} \rightarrow \mathbf{A}$ that holds over relationship schema **R**, where A is a non-prime attribute

(i.e., A is not an attribute in any candidate key)

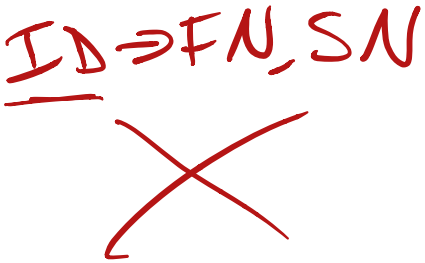
1. either $\mathbf{A} \in \mathbf{X}$ (it is trivial), or
2. **X** is a superkey for **R**, or
3. **X** is transitively dependent on a super key R

Easier to think of the opposite: There cannot be $\mathbf{X} \rightarrow \mathbf{A}$ where X is a partial candidate key for R

- Says nothing about non-prime to non-prime dependencies!

2NF Violations

ID	First Name	Surname	Telephone Number
123	Pooja	Singh	555-861-2025
123	Pooja	Singh	192-122-1111
456	San	Zhang	182-929-2929
456	San	Zhang	(555) 403-1659 Ext. 53
789	John	Zhang	555-808-9633



Tournament	Year	Winner	Winner's Birthplace
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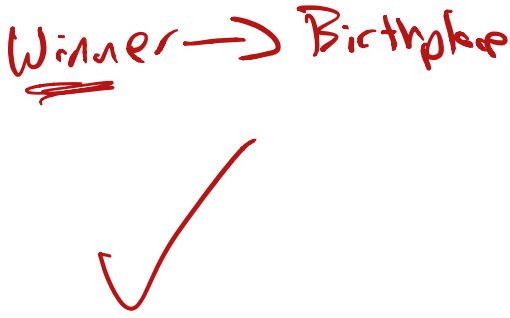

2NF Violations

ID	First Name	Surname	Telephone Number
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456	San	Zhang	(555) 403-1659 Ext. 53
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ID -> {First Name, LastName} Violates 2NF since **ID** is a partial Candidate Key

Tournament	Year	Winner	Winner's Birthplace
Indiana Invitational	1998	Al Fredrickson	Ohio
Cleveland Open	1999	Bob Albertson	New York
Des Moines Masters	1999	Al Fredrickson	Ohio
Indiana Invitational	1999	Chip Masterson	Kentucky

No 2NF violation

Redefining 3NF

Third Normal Form (3NF): For every **X** → **A** that holds over relationship schema **R**,

1. either **A** ∈ **X** (it is trivial), or
2. **X** is a superkey for **R**, or
3. **A** is a member of some key for **R**

Easier to think of the opposite: There cannot be **X** → **A** where X is not a full candidate key for R

"Every non-key attribute must provide a fact about the Key, the whole Key, and nothing but the Key… so help me Codd"

3NF Violations

Birthday → Age

Customer ID	First Name	Surname	Birthday	Age	Fav Color
123	Pooja	Singh	1/4/1984	37	Blue
456	San	Zhang	3/15/2001	19	Blue
789	John	Zhang	11/12/2006	14	Buff

Winner → Birthplace

Tournament	Year	Winner	Winner's Birthplace
Indiana Invitational	1998	Al Fredrickson	Ohio
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3NF Violations

Customer ID	First Name	Surname	Birthday	Age	Fav Color
123	Pooja	Singh	1/4/1984	37	Blue
456	San	Zhang	3/15/2001	19	Blue
789	John	Zhang	11/12/2006	14	Buff

Birthday->Age holds, but **Birthday** is not a superkey

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Indiana Invitational	1998	Al Fredrickson	Ohio
Cleveland Open	1999	Bob Albertson	New York
Des Moines Masters	1999	Al Fredrickson	Ohio
Indiana Invitational	1999	Chip Masterson	Kentucky

Winner -> Birthplace holds, but **Winner** is not a superkey

Normal Forms 1-3

1NF: Attributes should be atomic and tables should have no repeating groups

- *Prevents messiness within a cell and repetition of rows*

2NF: There cannot be **X** → **A** where X is a partial candidate key for R

- Doesn't forbid non-prime to non-prime dependencies
- *Prevents repetition of cells across rows*

3NF: There cannot be **X** → **A** where X is not a full candidate key for R

- Only allows dependencies on Keys
- *Prevents repetition of data within a row*